

## **A Fast Algorithm for the Cyber 205 to Simulate the 3D Ising Model**

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We describe a computer program that performs the Metropolis algorithm for the 3D Ising model at a peak speed of 98 million spin updates per second on a 2-pipe CDC Cyber 205. This speed is achieved using the special vector capabilities of the Cyber 205 and multispin coding techniques.

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**KEY WORDS:** Ising model, Monte Carlo method, multispin coding, vector computer.

### **INTRODUCTION**

This paper describes a new way to implement the Metropolis et al.<sup>(1)</sup> algorithm for Monte Carlo simulations of statistical systems with a few discrete degrees of freedom per variable. We describe this algorithm as implemented on a CYBER 205 and for the 3D Ising model. However, as should become obvious, our method can also be used on other computers and for other systems.

Each spin update in the Metropolis algorithm involves comparing the exponential of the change in the action with a random number. Since generating a random number takes about 20 ns on the Cyber 205, the best that a normal implementation of the algorithm can achieve is 50 million updates per second. We describe a method that avoids the bottleneck of making a floating point comparison with a random number. (Another method to avoid this bottleneck was proposed in Ref. 2.) Our method, in brief, works as follows:

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1. In our simulation, we use a variation of the multispin coding method<sup>(3)</sup> in which a single word contains one spin from  $n$  different systems. In our case,  $n$  is the word length which for the CYBER 205,  $n = 64$ . Hence we are simultaneously working on 64 different lattices.
2. We code the necessary information about each random number into two bits.
3. Each bit pair is used exactly once for each of the 64 lattices.
4. We use only logical commands for the update. Thus, 64 spins are updated together.

This method gives an algorithm speed of 98 megaflops in a 2-pipe Cyber 205. The fastest implementation for the 3D Ising model we are aware of does 218 megaflops on a DAP computer. However, this implementation is for the special case of a  $128 \times 128 \times 144$  lattice.<sup>(4)</sup>

Because of inherent Cyber 205 limitations on vector length, our current algorithm can only run on lattices of size up to  $50^3$ . To modify the code for larger lattices is easy. It involves slicing the lattice up into two dimensional planes and making the vector length equal half the number of spins in a plane. The algorithm currently achieves a speed of over 90 million spin updates per second for lattices of size greater than  $14^3$  on a 2-pipe Cyber 205. Our peak speed of 98 megaflops is achieved on a  $20^3$  lattice. We have used this program recently to reanalyze finite-size scaling for the 3D Ising model and compute the critical exponent  $\gamma$  to a few parts in a thousand.<sup>(5)</sup>

## THE ALGORITHM

The Ising model is defined as a collection of spins on the sites of the lattice. The action (energy) of the system is given by

$$S(\{s\}) = - \sum_{i,\hat{\mu}} s_i s_{i+\hat{\mu}} \quad (s = \pm 1) \quad (1a)$$

The spins  $s_i$  can have two possible values,  $\pm 1$ . The aim is to generate configurations of spins with joint probability distribution

$$Z = \sum_{\{s\}} e^{-\beta S(\{s\})} \quad (1b)$$

It is more convenient to store one spin per bit by using, in place of the variables  $s$ , the variables  $\sigma = (1 - s)/2$ , which take values 0 or 1. As a function of these variables, the action is

$$S(\{\sigma\}) = - \sum_{i,\hat{\mu}} (1 - 2x_{i,\hat{\mu}}) \quad (x_{i,\hat{\mu}} = 0, 1) \quad (2a)$$

with

$$x_{i,\mu} = XOR(\sigma_i, \sigma_{i+\mu}) \quad (\sigma = 0, 1) \quad (2b)$$

The change in the action on flipping the spin at site  $i$  is

$$\Delta S = S_{\text{final}} - S_{\text{initial}} = 12 - 4 \sum_{\pm\mu} x_{i,\mu} \quad (3)$$

where the sum goes over all neighbors of  $s_i$ . The Metropolis algorithm consists of the following

1. If  $\Delta S$  is nonpositive the spin flip should be accepted.
2. Otherwise, it should be accepted with probability  $e^{-\beta\Delta S}$ .

We implement this by logical commands as follows:

Define three bit variables  $B3V$ ,  $B2V$ ,  $B1V$  initialized to 0, 0, 1, respectively. Add the six values of  $x_{i,\mu}$  to the bits  $B3V$ ,  $B2V$ ,  $B1V$ , thinking of them as the third, second, and first bits of an integer. Note that this can be done with logical instructions since each  $x_{i,\mu}$  is either 0 or 1. The seven final values of the set  $B3V$ ,  $B2V$ ,  $B1V$  are shown in Table I. Notice that when the spin flip is to be accepted  $B3V=1$ . Using  $B3V$  as the acceptance criterion implements the first part of the Metropolis algorithm. If  $B3V$  is still zero after the additions, then we are dealing with the cases where  $\sum_{\pm\mu} x_{i,\mu}$  is 0, 1, or 2 and in this case, the spin should flip with probability  $e^{-\beta\Delta S}$ . To do this, we define two additional bits  $D2V$ ,  $D1V$ , called *demons* (*demon* variables similar to ours were first used by Creutz in the context of the microcanonical ensemble.<sup>(6)</sup>) which take the values (0, 1), (1, 0), and (1, 1) with probabilities  $p_{01}$ ,  $p_{10}$ , and  $p_{11}$  given by

$$\begin{aligned} p_{01} &= e^{-4\beta} - e^{-8\beta} \\ p_{10} &= e^{-8\beta} - e^{-12\beta} \\ p_{11} &= e^{-12\beta} \end{aligned} \quad (4)$$

Table I. Logical Implementation of the Metropolis Algorithm

$\Delta S$	$\sum_{\pm\mu} x_{i,\mu}$	$B3$	$B2$	$B1$
		0	0	0
12	0	0	0	1
8	1	0	1	0
4	2	0	1	1
0	3	1	0	0
-4	4	1	0	1
-4	5	1	1	0
-12	6	1	1	1

The integer formed with  $D1V$  and  $D2V$  as the first and second bit is now added to the integer formed by  $B3V$ ,  $B2V$ , and  $B1V$ . Once again, if  $B3V$  is unity, the flip is accepted. Notice that this will happen with probability  $e^{-4\beta}$ ,  $e^{-8\beta}$ , and  $e^{-12\beta}$  when  $\Delta S$  is 4, 8, or 12, respectively. This implements the second part of the Metropolis algorithm.

Now we discuss the implementation of this algorithm on the Cyber 205. (The code is shown in Fig. 2.) As mentioned earlier, our program updates 64 different lattices simultaneously. Spins occupying the same coordinates on each of the 64 lattices are stored in the same word. The spins are labeled even/odd in a checkerboard pattern and updated in two steps: first all even spins are updated and then all the odd ones. The vector length of the pipelined arrays is half the total number of spins on a lattice. The variables  $B1V$ ,  $B2V$ , and  $B3V$ , as well as the demons  $D1V$  and  $D2V$ , are also arrays of the same length. Clearly, the algorithm can be pipelined in the Cyber 205 if one can arrange the arrays  $D1V$  and  $D2V$  such that their entries (in pairs), take on the values (0, 1), (1, 0), and (1, 1) with probabilities given by eq. 4.

The crucial part of the algorithm is to get such a distribution of demon bits. This is done in subroutine DEMETRO. First, 64 vectors of random numbers are generated one after the other using a shift register random number generator.<sup>(7)</sup> Each of these vectors is used to define one string of demon pairs ( $D1V$ ,  $D2V$ ) distributed according to eq. 4. The final vectors for  $D1V$  and  $D2V$  are obtained by merging together, in the 64-bit positions of the words, the 64 strings thus obtained. After each half sweep (update of the odd or even sites of all the 64 lattices), the demon pairs ( $D1V$ ,  $D2V$ ) are subjected to a random GATHER and a random shift. The random shifts are arranged so that each string of demons is used in each of the lattices exactly once but in a different order for each lattice. After 64 half sweeps, the demons are reinitialized. This means that we use one random number once for each of the 64 lattices.

In *any* Metropolis run, one uses a finite sequence of random numbers to compare to  $e^{-\beta\Delta S}$ . In such a finite run, the set of random numbers used is not uniformly distributed in the interval (0, 1) and this means that the  $\beta$  value is not the desired one but rather, some other value  $\beta_{\text{eff}}$ . For our simulation, these effective values of  $\beta$  are given by

$$\begin{aligned} n_1/n &= e^{-4\beta_{\text{eff},1}} \\ n_2/n &= e^{-8\beta_{\text{eff},2}} \\ n_3/n &= e^{-12\beta_{\text{eff},3}} \end{aligned} \quad (5)$$

where  $n$  is the total number of random numbers used and  $n_1$ ,  $n_2$ , and  $n_3$  are

the number of times these random numbers are smaller than  $e^{-4\beta}$ ,  $e^{-8\beta}$ , and  $e^{-12\beta}$ , respectively.

This expected shift in  $\beta$  can be easily understood theoretically. Given  $n$  random numbers in the interval  $(0, 1)$ , the probability for  $n_i$  of them being in an interval of length  $p_k = e^{-4k\beta}$  is given by the binomial distribution

$$P_k(n_i) = \binom{n}{n_i} p_k^{n_i} (1 - p_k)^{n - n_i} \tag{6}$$

If we define  $x \equiv n_i/n$ , then this probability distribution has mean  $\bar{x} = p_k$  and standard deviation  $\sigma_k = \sqrt{p_k(1 - p_k)/n}$ . Note that the error in  $\beta$  is statistical and proportional to  $1/\sqrt{n}$  and disappears in the limit of an infinite simulation. It is easy to correct for this shift in  $\beta$ . This is done by correcting each order parameter  $O$  according to

$$O(\beta) = O(\beta_{\text{eff}}) + \left. \frac{\partial O}{\partial \beta} \right|_{\beta_{\text{eff}}} (\beta - \beta_{\text{eff}}) \tag{7}$$

The reason we need to worry about this error (which is actually present even in standard Metropolis updating but is less relevant there) is the following: In our simulation, each of the 64 lattices uses the same set of random numbers. Therefore, each of the lattices has the same  $\beta_{\text{eff}}$  on average and the effects of the shift in  $\beta$  add coherently. In a standard Metropolis run on 64 lattices using different random number streams on each, the effective length of the random number stream is 64 times our length. The shift in  $\beta$  then is smaller by a factor of 8.

We obtained one single value of the effective  $\beta$  using the average demon action  $E_{\text{demon}}$  via the formula

$$\begin{aligned} & \frac{4e^{-4\beta_{\text{eff},1}} + 8e^{-8\beta_{\text{eff},2}} + 12e^{-12\beta_{\text{eff},3}}}{1 + e^{-4\beta_{\text{eff},1}} + e^{-8\beta_{\text{eff},2}} + e^{-12\beta_{\text{eff},3}}} = \langle E_{\text{demon}} \rangle \\ & = \frac{4e^{-4\beta_{\text{eff}}} + 8e^{-8\beta_{\text{eff}}} + 12e^{-12\beta_{\text{eff}}}}{1 + e^{-4\beta_{\text{eff}}} + e^{-8\beta_{\text{eff}}} + e^{-12\beta_{\text{eff}}}} \end{aligned} \tag{8}$$

We have tested these ideas on the two-dimensional Ising model where exact answers are known. This is shown in Fig. 1. The 40 points in the graphs correspond to different runs, each of them coming from 20,000 sweeps in a set of 64 lattices of size  $20^2$  at  $\beta = 0.44$  and averaging over the last 10,000 sweeps. In Fig. 1 we show average values of  $\langle S^2 \rangle$  before and after the correction. The errors were obtained by using the 64 average values from the 64 independent lattices. The uncorrected data values (Fig. 1a), as expected, are scattered around a shifted  $\beta$  with a Gaussian

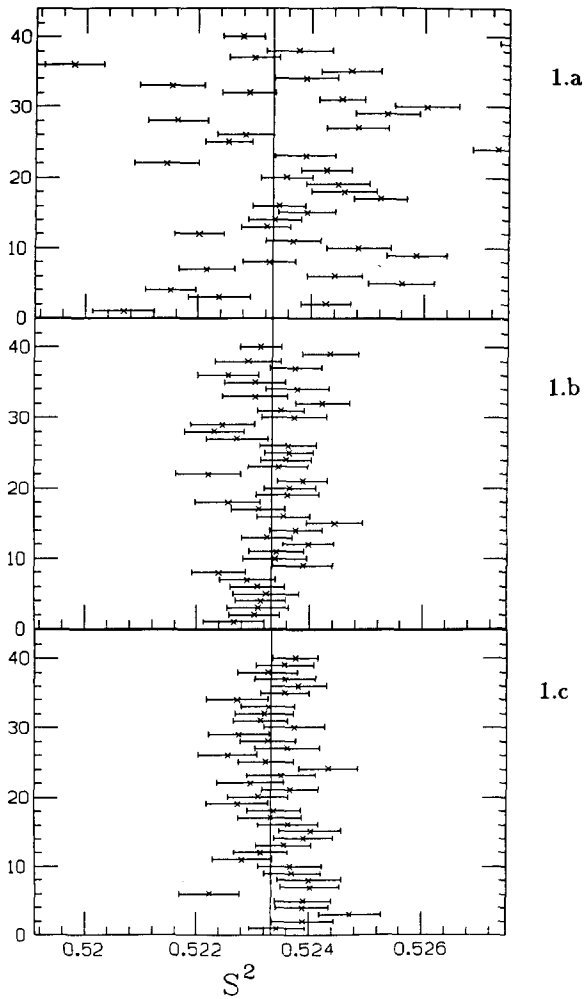


Fig. 1. 40 data points for  $\langle S^2 \rangle$  on  $20^2$  lattices at  $\beta = .44$ . Each point corresponds to averaging for 10,000 sweeps on the set of 64 lattices after thermalizing for 10,000 sweeps. The error bars were obtained from fluctuations between the 64 data sets from the different lattices. (a) Shows the data without correction; (b) Shows the data after correcting for the shift in  $\beta$  and using eq. 10. (c) Shows the data corresponding to the same experiment but using a different random number stream for all 64 lattices. (a conventional, but slower, standard Metropolis simulation). The vertical line that crosses the graphs corresponds to the exact value. Note that the error bars are all of almost the same size.

```

c 3-d ising model program
c
c this program performs the metropolis algorithm on cubic lattices
c of size up to 50*3, it uses periodic boundary conditions.
c the lattice size must be even.

      program ising ( input , output , tape6 , tape7 )

c define parameters
c
c      l      = lattice size (must also be def.in sub. metro)
c      init   = initialization parameter (0:ordered start,
c              1:disordered start,6:reads lattice from tape6)
c      istor  = (1,0):(does,does not) write last conf. to tape7
c      iseed  = seed for the r.n.g. (if =0 ==> default seed.)
c      nb     = # of beta (= 1/Temperature) values processed
c      beta   = initial beta
c      dibe   = successive decrements in beta.
c      nsba   = # of sweeps to thermalize
c      nsw    = # of sweeps during averaging
c      nav    = number of sweeps between measurements.
c      np,nq  = prim. trinomial for the shift register random
c              number generator (must also be defined in the sub. metro)
c              ( nh > np > nq )
c
      parameter ( l      = 20
+               ,init   = 1
+               ,istor  = 0
+               ,iseed  = 7893789324783
+               ,nb     = 1
+               ,beta   = .22165
+               ,dibe   = .0
+               ,nsba   = 25000
+               ,nsw    = 25000
+               ,nav    = 100
+               ,np     = 2281
+               ,nq     = 715
+               , ns=1**3, lm=1-1, nh=ns/2, nhp=nh+1, nt=nh+np)

c VARIABLES IN THE PROGRAM
c bx,by,bz: Indexing vectors to implement the periodic boundary conditions
c xrand: vector for the shift register random number generator
c xv,xv1,xv2,dn1v,dn2v: help vectors (to store intermediate results)
c spin: vector of ising spin variables for the 64 lattices.
c b1v,b2v,b3v: vectors used in the updating
c d1v,d2v: demon vectors
c nr,nrp,num,nr1,nr2: vectors used in demon permutations.
c ran: vector for congruential random number generator.
c tf,ti,tm,ta,timet,t64,tt: scalars for timing.
c i64,n64: vectors used to permute demons d1v,d2v.
c am,s,r,abe,ntb: variables used in measurements.
c cv,cv1,cv2: bit vectors used to update the demons

      integer bx,by,bz,xrand,xv,xv1,xv2,spin,b1v,b2v,b3v
c      ,d1v,d2v,dn1v,dn2v
c      ,dx,dy,dz,x,x1,x2,dxra,dsp1,dsp2,b1,b2,b3,d1,d2
c      ,dn1,dn2,dnr,dnpr,dnr1,dnr2

      descriptor dx,dy,dz,x,x1,x2,dxra,dsp1,dsp2,b1,b2,b3,d1,d2
c      ,dn1,dn2,dnr,dnpr,dnr1,dnr2,ia,ib,ic,isd

      bit cv,cv1,cv2

      common/blk/ spin(ns),bx(nh),by(nh),bz(nh),xv(nh),xv1(nh),xv2(nh)
c      ,b1v(nh),b2v(nh),b3v(nh),d1v(nh),d2v(nh),dn1v(nh),dn2v(nh)
c      ,nrp(nh),num(nh),nr2(nh),nr1(nh),nr(nh),ran(nh),tf,ti,tm,ta
c      ,timet,t64,i64(64),n64(64),am(64),s(64),r(16,64),abe(3)
c      ,ntb,xrand(nt),cv(2*nh),cv1(2*nh),cv2(2*nh)

      data tf,ti,tm,ta,timet,tt,t64,(abe(i),i=1,3),ntb /10*.0,0/

c ASSIGN DESCRIPTORS TO ARRAYS:

      data dsp1,dsp2 / spin(1:nh),spin(nhp:nh) /
      data dx,dy,dz / bx(1:nh),by(1:nh),bz(1:nh) /
      data b1,b2,b3 / b1v(1:nh),b2v(1:nh),b3v(1:nh) /
      data d1,d2 / d1v(1:nh),d2v(1:nh) /
      data dnr,dnpr,dnr1,dnr2 / nr(1:nh),nrp(1:nh),nr1(1:nh),nr2(1:nh)/
      data dn1,dn2,x,x1,x2 / dn1v(1:nh),dn2v(1:nh),xv(1:nh),xv1(1:nh)
+      , x2v(1:nh) /

```

Fig. 2. Listing of the code.

```

assign ia , xrand( 1 : nh )
assign ib , xrand( nq+1 : nh )
assign ic , xrand( nh+1 : np )
assign isd , xrand( 1 : np )
assign dxra , xrand( np+1 : nh )

tt1=second()
74 do 74 i=1,nh
   num(i)=i
c initialization of: random number seeds, the lattice, and the indexing
c vectors

   if(iseed.ne.0)call ranset(iseed)
   call raninit(xrand,np)
   call initlat(init,betao,nso,dsp1,dsp2,ran,x1,spin,nh,ns,l)
   call index(bx,by,bz,l)
c begin simulation

   do 1234 iib = 1,nb

   b = beta - dibe*iib
   if ( b.lt.0. ) b=-b

   call metro( nsba , 0 , b ,dsp1,dsp2,dx,dy,dz,x,x1,x2,b1,b2,b3
+ ,d1,d2,dn1,dn2,dnrd,dnr1,dnr2,dnr,dxra,ia,ib,ic,isd)
   call metro( nsw , nav , b ,dsp1,dsp2,dx,dy,dz,x,x1,x2,b1,b2,b3
+ ,d1,d2,dn1,dn2,dnrd,dnr1,dnr2,dnr,dxra,ia,ib,ic,isd)
c write the data to the output file

   if(nsw.ne.0)call data(b,nsw,nsba,nav,init,betao,nso,iib,ntb
+ ,abe,ns,r)

1234 continue
c simulation ends. Stores the last lattice

   if(istar.eq.1)then
ns=nsba+nsw
   if(dibe.eq.0)then
nsn=nb*nsn
   if(betao.eq.beta)nsn=nso+nsn
   endif
   write(7,*)l,beta,nsn
   do 77 i=1,ns
77 write(7,177) spin(i)
177 format(z16)
   endif
c timing information:

   tt=second()-tt1
   print*,'time spent in the main parts of the program :
   print*
   print565,' program ising = ',tt
   print565,' in updating = ',timet
   print565,' in measuring = ',ta
   print565,' sub. demetro = ',tf
   print565,' sub. initnr = ',ti
   print565,' sub. init64 = ',t64
   print565,' sub. meas = ',tm
   print*
   print566,(64e-6*ns*nb*(nsba+nsw))/timet
565 format('x,a20,f10.2,' seconds')
566 format(' running at ',f5.1,' megaflips')
   print*
   stop
   end

   subroutine index(bx,by,bz,l)
   integer bx(*),by(*),bz(*)
c defines the index vectors to implement periodic boundary conditions.

   lm = l - 1

```

Fig. 2 (continued)



```

do 1 iz = 0,lm
izm = mod( iz+lm , l )
do 1 iy = 0,lm
iym = mod( iy+lm , l )
do 1 ix = 0,lm
ixn = mod( ix+1 , l )
if ( mod(iy+iz,2) .eq. 0 )ixn = mod( ix+lm , l )

if(mod(ix+iy+iz,2).eq.0)then
n = ( iz *1 + iy )*(1/2) + (ix -mod(ix ,2) )/2 + 1
bx(n) = ( iz *1 + iy )*(1/2) + (ixn-mod(ixn,2) )/2
by(n) = ( iz *1 + iym)*(1/2) + (ix -mod(ix ,2) )/2
bz(n) = ( izm*1 + iy )*(1/2) + (ix -mod(ix ,2) )/2
endif

1 continue
return
end

subroutine initnr(nr,ti,nrp,num,ran,nh)

c creates a vector (nr) containing a random permutation of
c the integers 0 to nh-1

dimension nr(*) ,nrp(*) ,num(*) ,ran(*)

ti1=second()

call vranf(ran,nh)
nrp(1:nh)=num(1:nh)*ran(1:nh)+1
call q8vrevv('x'08',,nrp(1:nh),...nr(1:nh))
nrp(1:nh)=num(1:nh)-1

do 1 i = 1,nh
n=nr(i)
nr(i) = nrp(n)
nv=nh-i-n+1
nrp(n;nh-i-n+1)=nrp(n+1;nh-i-n+1)
1 continue

ti=ti+second()-ti1
return
end

subroutine metro( nsw , nav , b ,dsp1,dsp2,dx,dy,dz,x,x1,x2,b1
+ ,b2,b3,d1,d2,dn1,dn2,dnr,dnr1,dnr2,dnr,dxra,ia,ib,ic,isd)
c metro( nsw , nav , b , ... ) does nsw sweeps measuring after each nav
c sweeps at beta = b. if nav = 0 no measurements are done. The other
c arguments are passive. They are descriptors needed by this routine.
parameter (l=20,ns=i*3,lm=i-1,nh=ns/2,nhp=nh+1)
parameter(np=2281,nq=715,nt=nh+np)
integer bx,by,bz,xrand,xv,x1v,x2v,spin,b1v,b2v,b3v
c ,d1v,d2v,dn1v,dn2v
c ,dx,dy,dz,x,x1,x2,dxra,dsp1,dsp2,b1,b2,b3,d1,d2
c ,dn1,dn2,dnr,dnrp,dnr1,dnr2
c descriptor dx,dy,dz,x,x1,x2,dxra,dsp1,dsp2,b1,b2,b3,d1,d2
c ,dn1,dn2,dnr,dnrp,dnr1,dnr2,ia,ib,ic,isd
bit cv,cv1,cv2
common/blk/ spin(ns),bx(nh),by(nh),bz(nh),xv(nh),x1v(nh),x2v(nh)
c ,b1v(nh),b2v(nh),b3v(nh),d1v(nh),d2v(nh),dn1v(nh),dn2v(nh)
c ,nrp(nh),num(nh),nr2(nh),nr1(nh),nr(nh),ran(nh),tf,ti,tm,ta
c ,timet,t64,i64(64),n64(64),am(64),s(64),r(16,64),abe(3),ntb
c ,xrand(nt),cv(2*nh),cv1(2*nh),cv2(2*nh)

c set counters

if( nav.ne.0 )then
naver = 0
r(1,1;1024)=.0
endif
is=64
isnr=63

c does nsw sweeps

do 99999 niter = 1,nsw
t2 = second()

```

Fig. 2 (continued)

```

c creates vectors of demons and random permutations
    if( is.eq.64 )then
    is=0
    call demetro(b,nav,d1,d2,ia,ib,ic,isd,dxra,xrand,tf,ntb
+
    ,d1v,d2v,abe,nh,np+1,cv,cv1,cv2)
    t641=second()
    call init64( i64 , n64 )
    t64=t64+second()-t641
    xv(1)=nr2(nh)
    xv(2:nh-1)=nr2(1:nh-1)
    dnr2=x
    call q8 vtovx (x' 00 ',, dnr2 ,, dnr1 ,, x )
    dnr1=x
    isnr=isnr+1
    if(isnr.eq.64)then
    call initnr( nr ,ti,nrp,num,ran,nh)
    call initnr( nr1,ti,nrp,num,ran,nh)
    call initnr( nr2,ti,nrp,num,ran,nh)
    isnr=0
    endif
endif

c updates even sites
c adds the six products x to [b3,b2,b1] = [0,0,1]

    call q8 xor v(x' 00 ',, dsp1 ,, dsp2 ,, x )
    call q8 vxtov (x' 00 ',, dy ,, dsp2 ,, x1 )
    call q8 xor v(x' 00 ',, x1 ,, dsp1 ,, x1 )
    call q8 and v(x' 01 ',, x1 ,, x ,, b2 )
    call q8 xor v(x' 00 ',, x1 ,, x ,, b1 )
    call q8 vxtov (x' 00 ',, dz ,, dsp2 ,, x )
    call q8 xor v(x' 00 ',, x ,, dsp1 ,, x )
    call q8 vtovx (x' 00 ',, dx ,, dsp2 ,, x1 )
    call q8 xor v(x' 00 ',, x1 ,, dsp1 ,, x1 )
    call q8 and v(x' 01 ',, x1 ,, x ,, dn2 )
    call q8 xor v(x' 00 ',, x1 ,, x ,, dn1 )
    call q8 and v(x' 01 ',, dn2 ,, b2 ,, b3 )
    call q8 and v(x' 01 ',, dn1 ,, b1 ,, x )
    call q8 xor v(x' 00 ',, dn2 ,, b2 ,, b2 )
    call q8 xor v(x' 00 ',, x ,, b2 ,, b2 )
    call q8 xor v(x' 00 ',, dn1 ,, b1 ,, b1 )
    call q8 vtovx (x' 00 ',, dy ,, dsp2 ,, x )
    call q8 xor v(x' 00 ',, x ,, dsp1 ,, x )
    call q8 vtovx (x' 00 ',, dz ,, dsp2 ,, x1 )
    call q8 xor v(x' 00 ',, x1 ,, dsp1 ,, x1 )
    call q8 or v(x' 02 ',, x1 ,, x ,, dn2 )
    call q8 xorn v(x' 07 ',, x1 ,, x ,, dn1 )
    call q8 and v(x' 01 ',, dn2 ,, b2 ,, x1 )
    call q8 xor v(x' 00 ',, x1 ,, b3 ,, b3 )
    call q8 xor v(x' 00 ',, dn2 ,, b2 ,, b2 )
    call q8 and v(x' 01 ',, dn1 ,, b1 ,, x1 )
    call q8 and v(x' 01 ',, x1 ,, b2 ,, x2 )
    call q8 xor v(x' 00 ',, x2 ,, b3 ,, b3 )
    call q8 xor v(x' 00 ',, x1 ,, b2 ,, b2 )
    call q8 xor v(x' 00 ',, dn1 ,, b1 ,, b1 )

c keeps third demon ( now b3=1 if inc(s) < 0 )
c gets a new set of demons

    call q8 vtovx (x' 00 ',, dnr1 ,, dnr ,, dnrp )
    is = is + 1
    ishi=i64(is)
    call q8 vtovx (x' 00 ',, dnrp ,, d1 ,, dn1 )
    call q8 shiftv(x' 08 ',, dn1 ,, ishi ,, d1 )
    call q8 vtovx (x' 00 ',, dnrp ,, d2 ,, dn2 )
    call q8 shiftv(x' 08 ',, dn2 ,, ishi ,, d2 )

c now checks what happens when inc(s) > 0 (that is when b3=0 )
c adds demons [d2,d1] to [x=0,b2,b1]

    call q8 and v(x' 01 ',, d2 ,, b2 ,, x )
    call q8 xor v(x' 00 ',, d2 ,, b2 ,, b2 )
    call q8 and v(x' 01 ',, d1 ,, b1 ,, x1 )
    call q8 and v(x' 01 ',, x1 ,, b2 ,, x2 )
    call q8 xor v(x' 00 ',, x2 ,, x ,, x )

c if now x = 1 also flips

```

Fig. 2 (continued)

```

c x is the acceptance criterion
      call q8 or v(x' 02 ',, x ,, b3 ,, x )

c updates the spins
      call q8 xor v(x' 00 ',, dsp1 ,, x ,, dsp1)

c updates odd sites
      call q8 xor v(x' 00 ',, dsp2 ,, dsp1 ,, x )
      call q8 vxtov(x' 00 ',, dy ,, dsp1 ,, x1 )
      call q8 xor v(x' 00 ',, x1 ,, dsp2 ,, x1 )
      call q8 and v(x' 01 ',, x1 ,, x ,, b2 )
      call q8 xor v(x' 00 ',, x1 ,, x ,, b1 )
      call q8 vxtov(x' 00 ',, dz ,, dsp1 ,, x )
      call q8 xor v(x' 00 ',, x ,, dsp2 ,, x )
      call q8 vxtov(x' 00 ',, dx ,, dsp1 ,, x1 )
      call q8 xor v(x' 00 ',, x1 ,, dsp2 ,, x1 )
      call q8 and v(x' 01 ',, x1 ,, x ,, dn2 )
      call q8 xor v(x' 00 ',, x1 ,, x ,, dn1 )
      call q8 and v(x' 01 ',, dn2 ,, b2 ,, b3 )
      call q8 and v(x' 01 ',, dn1 ,, b1 ,, x )
      call q8 xor v(x' 00 ',, dn2 ,, b2 ,, b2 )
      call q8 xor v(x' 00 ',, x ,, b2 ,, b2 )
      call q8 xor v(x' 00 ',, dn1 ,, b1 ,, b1 )
      call q8 vtovx(x' 00 ',, dy ,, dsp1 ,, x )
      call q8 xor v(x' 00 ',, x ,, dsp2 ,, x )
      call q8 vtovx(x' 00 ',, dz ,, dsp1 ,, x1 )
      call q8 xor v(x' 00 ',, x1 ,, dsp2 ,, x1 )
      call q8 or v(x' 02 ',, x1 ,, x ,, dn2 )
      call q8 xorn v(x' 07 ',, x1 ,, x ,, dn1 )
      call q8 and v(x' 01 ',, dn2 ,, b2 ,, x1 )
      call q8 xor v(x' 00 ',, x1 ,, b3 ,, b3 )
      call q8 xor v(x' 00 ',, dn2 ,, b2 ,, b2 )
      call q8 and v(x' 01 ',, dn1 ,, b1 ,, x1 )
      call q8 and v(x' 01 ',, x1 ,, b2 ,, x2 )
      call q8 xor v(x' 00 ',, x2 ,, b3 ,, b3 )
      call q8 xor v(x' 00 ',, x1 ,, b2 ,, b2 )
      call q8 xor v(x' 00 ',, dn1 ,, b1 ,, b1 )
      call q8 vtovx(x' 00 ',, dnr1 ,, dnrp ,, dnr )
      is=is+1
      ishi=i64(is)
      call q8 vtovx(x' 00 ',, dnr ,, d1 ,, dn1 )
      call q8 shiftv(x' 08 ',, dn1 ,, ishi ,, d1 )
      call q8 vtovx(x' 00 ',, dnr ,, d2 ,, dn2 )
      call q8 shiftv(x' 08 ',, dn2 ,, ishi ,, d2 )
      call q8 and v(x' 01 ',, d2 ,, b2 ,, x )
      call q8 xor v(x' 00 ',, d2 ,, b2 ,, b2 )
      call q8 and v(x' 01 ',, d1 ,, b1 ,, x1 )
      call q8 and v(x' 01 ',, x1 ,, b2 ,, x2 )
      call q8 xor v(x' 00 ',, x2 ,, x ,, x )
      call q8 or v(x' 02 ',, x ,, b3 ,, x )
      call q8 xor v(x' 00 ',, dsp2 ,, x ,, dsp2)

c times the updating
      timet = timet + second()-t2

c takes averages
      if( nav.eq.0 )goto 99999
      if( mod(niter,nav).ne.0 )goto 99999
      ta3=second()
      call meas(dsp1,dsp2,dx,dy,dz,x,x1,x2,am,s,nh,tm)
      naver = naver+1
      do 771 i=1,64
      ss=s(i)
      rm=abs(am(i))
      rm2=rm*rm
      rm4=rm2*rm2
      s2=ss*ss
      s4=s2*s2
      r( 1,i) = r( 1,i) + am(i)
      r( 2,i) = r( 2,i) + rm
      r( 3,i) = r( 3,i) + rm2
      r( 4,i) = r( 4,i) + rm4
      r( 5,i) = r( 5,i) + ss
  
```

Fig. 2 (continued)

```

r( 6,i) = r( 6,i) + s2
r( 7,i) = r( 7,i) + s4
r( 8,i) = r( 8,i) + rm*ss
r( 9,i) = r( 9,i) + rm2*ss
r(10,i) = r(10,i) + rm4*ss
771 r(11,i) = r(11,i) + ss*s2
r(12,i) = r(12,i) + ss*s4

ta=ta+second()-ta3

99999 continue

if( nav.gt.nsw .or. nav.eq.0 )return
r(1,1;1024) = r(1,1;1024)/naver
return
end

subroutine demetra(b,nav,d1,d2,ia,ib,ic,isd,dxra,xrand,tf,ntb
+ ,d1v,d2v,abe,nh,npm,cv,cv1,cv2)

c sets the demons with the right probabilities

integer xrand(*),d1v(*),d2v(*),dxra,d1,d2
descriptor dxra,d1,d2,ia,ib,ic,isd
bit cv(*),cv1(*),cv2(*)
dimension abe(3),be(3)
half precision ma(2),e1,e2,e3
equivalence(mar,ma(1))

tf1=second()
nh2=2*nh
d1 = 0
d2 = 0
be(1;3)=0.
rrr=1./nh2

c gets normalized boltzman factors

e1 = 2.**23 *exp( -4.*b )
e2 = 2.**23 *exp( -8.*b )
e3 = 2.**23 *exp( -12.*b )

c loops over the 32 bits of the halfwords (uses half precision)

do 1 i=0,31
mar=shift(1,i)

c gets random numbers using a shif register random number generator

call q8 xor v(x'00' ,, ia ,, ib ,, dxra )
call q8 vto v(x'00' ,, ic ,,, isd )

c sets the demons and measures the effective beta

call q8cmlt(x'88' ,, xrand(npm;nh2) ,,e2 ,cv(1;nh2))
call q8cmlt(x'88' ,, xrand(npm;nh2) ,,e1 ,cv1(1;nh2))
call q8cmlt(x'88' ,, xrand(npm;nh2) ,,e3 ,cv2(1;nh2))

if(nav.ne.0)then
be(1)=be(1)+q8sclt(cv1(1;nh2))*rrr
be(2)=be(2)+q8sclt(cv(1;nh2))*rrr
be(3)=be(3)+q8sclt(cv2(1;nh2))*rrr
endif

call q8xor v(x'88' ,, d2v(1;nh2) ,,ma(2),cv(1;nh2),d2v(1;nh2))
call q8andn(x'00' ,, cv1(1;nh2) ,, cv2(1;nh2) ,, cv1(1;nh2))
call q8xor(x'00' ,, cv1(1;nh2) ,, cv(1;nh2) ,, cv(1;nh2))
call q8xor v(x'88' ,, d1v(1;nh2) ,,ma(2),cv(1;nh2),d1v(1;nh2))

1 continue

if(nav.ne.0)then
abe(1;3)=abe(1;3)+be(1;3)/32.
ntb=ntb+1
endif
tf=tf+second()-tf1
return
end

```

Fig. 2 (continued)

```

subroutine meas(dsp1,dsp2,dx,dy,dz,x,x1,x2,am,s,nh,tm)
c measures the magnetization and the action for the 64 lattices.

integer dx,dy,dz,x,x1,x2,dsp1,dsp2
descriptor dx,dy,dz,x,x1,x2,dsp1,dsp2
dimension am(*),s(*)

tm1=second()

do 2 i = 1,64
ma = shift( 1 , i-1 )
ls = mod( 65-i , 64 )

c measures the magnetization

call q8 and v(x' 09 ' , , dsp1 , , ma , , x1 )
call q8 shiftv(x' 08 ' , , x1 , , ls , , x1 )
am(i) = q8ssum( x1 )
call q8 and v(x' 09 ' , , dsp2 , , ma , , x2 )
call q8 shiftv(x' 08 ' , , x2 , , ls , , x2 )
am(i) = am(i) + q8ssum( x2 )

c measures the action

call q8 xor v(x' 00 ' , , x1 , , x2 , , x )
s(i) = q8ssum( x )
call q8 vxtov (x' 00 ' , , dy , , x2 , , x )
call q8 xor v(x' 00 ' , , x , , x1 , , x )
s(i) = s(i) + q8ssum( x )
call q8 vxtov (x' 00 ' , , dz , , x2 , , x )
call q8 xor v(x' 00 ' , , x , , x1 , , x )
s(i) = s(i) + q8ssum( x )
call q8 vtovx (x' 00 ' , , dx , , x2 , , x )
call q8 xor v(x' 00 ' , , x , , x1 , , x )
s(i) = s(i) + q8ssum( x )
call q8 vtovx (x' 00 ' , , dy , , x2 , , x )
call q8 xor v(x' 00 ' , , x , , x1 , , x )
s(i) = s(i) + q8ssum( x )
call q8 vtovx (x' 00 ' , , dz , , x2 , , x )
call q8 xor v(x' 00 ' , , x , , x1 , , x )
s(i) = s(i) + q8ssum( x )

2 continue

am(1:64) = 1. - am(1:64)/nh
s(1:64) = 1. - s(1:64)/(3.*nh)
tm=tm+second()-tm1
return
end

subroutine data(b,nsw,nsba,nav,init,betao,nso,iib,ntb,abe,ns,r)

character a(130)
dimension t(16),v(16),r(16,64),abe(3)
data a/'130*'-'/'

c writes the header

print*,a
print * , 'beta = ',b
if(iib.eq.1)then
if(init.eq.0)print*, 'initial lattice ordered'
if(init.eq.1)print*, 'initial lattice disordered'
if(init.eq.6)then
print*, 'initial lattice read from tape 6, which was thermalized'
+ , ' during ',nso,' sweeps at beta = ',betao
endif
else
print*, 'initial lattice from previous run'
endif
print*, 'thermalizing during ',nsba,' sweeps'
print*, 'taking averages after every ',nav,' sweeps during '
+ ',nsw,' sweeps'
print*,a

c calculates the data for every lattice

```

Fig. 2 (continued)

```

t(1:16) = .0
v(1:16) = .0
do i = 1,64
  r(13,i) = ns * (r(3,i) - r(2,i)*r(2,i))
  r(14,i) = r(3,i) * ns
  r(15,i) = r(4,i)/(r(3,i)*r(3,i))-3.
  r(16,i) = 3.*ns * (r(6,i) - r(5,i)*r(5,i))
  t(1:16) = t(1:16) + r(1,i:16)
  v(1:16) = v(1:16) + r(1,i:16)*r(1,i:16)
8   print 8 ,i,r(1,i),r(2,i),r(5,i),r(13,i),r(14,i),r(16,i),r(15,i)
+   format(' lat ',i2,' m= ',f10.7,' |m|= ',f10.7,' s= ',f10.7
1   + , ' x= ',f11.6,' xp= ',f11.6,' c= ',f11.6,' gr= ',f11.6)
  continue

c writes the data averaged from the 64 lattices

do 53 i=1,16
  t(i) = t(i) /64.
53  v(i) = sqrt( v(i)/64. - t(i)*t(i) ) /8.
  print*,o
  print * , 'data from averaging the 64 lattices :
  nf=0
888  print *
  print11,' m = ',t(1),' |m| = ',t(2),' m+2 = ',t(3),' m+4 = ',t(4)
  print12,' +/- ',v(2),' +/- ',v(3),' +/- ',v(4)
  print *
  print12,' s = ',t(5),' s+2 = ',t(6),' s+4 = ',t(7)
  print12,' +/- ',v(5),' +/- ',v(6),' +/- ',v(7)
  print *
  print13,' xp = ',t(14),' x = ',t(13),' c = ',t(16),' gr = ',t(15)
  print13,' +/- ',v(14),' +/- ',v(13),' +/- ',v(16),' +/- ',v(15)
  if(nf.eq.1)goto 889
  print *
11  format( 5x,4(10x,a6,f11.8))
12  format(32x,3(10x,a6,f11.8))
13  format(15x,2(a5,f12.5,10x),a5,f12.6,10x,a5,f12.7)

c corrects the data for shifted value of beta.

abe(1:3)=abe(1:3)/ntb
coe=(abe(1)+2.*abe(2)+3.*abe(3))/(1.+tabe(1)+tabe(2)+tabe(3))
z1=0.
z2=1.
do 75 i=1,100
  z=(z1+z2)*.5
  ed=(z+2.*z*z+3.*z*z*z)/(1.+z+z*z+z*z*z)
  if(ed.eq.coe)goto 777
  if(ed.lt.coe)then
    z1=z
  else
    z2=z
  endif
75  continue

777  ba=(-1./4.)*log(z)
  rib=3.*ns*(b-ba)
  t6=t(6)
  t(2) = t(2) + (t(8) - t(2)*t(5)) * rib
  t(3) = t(3) + (t(9) - t(3)*t(5)) * rib
  t(4) = t(4) + (t(10) - t(4)*t(5)) * rib
  t(6) = t(6) + (t(11) - t(6)*t(5)) * rib
  t(7) = t(7) + (t(12) - t(7)*t(5)) * rib
  t(5) = t(5) + (t6 - t(5)*t(5)) * rib
  t(14) = t(3) * ns
  t(15) = t(4)/(t(3)*t(3))-3.
  t(13) = ns * (t(3) - t(2)*t(2))
  t(16) = 3.*ns * (t(6) - t(5)*t(5))

  print * , 'effective beta = ', ba
  print *
  print * , 'corrected data:'
  nfb=1
  goto 888
889  print*,o
  print*
  ntb=0
  abe(1:3)=0.

```

Fig. 2 (continued)

```

return
end

subroutine initlat(init,betao,nso,dsp1,dsp2,ran,x1,spin,nh,ns,l)
c initializes the lattices

integer spin,x1,dsp1,dsp2
descriptor x1,dsp1,dsp2
dimension spin(*),ran(*)

dsp1 = 0
dsp2 = 0
if ( init.eq. 1 )then

do 1 i = 1,64
ls=mod(65-i,64)
x1=0
call vranf(ran,nh)
where (ran(1:nh).ge..5)x1=1
call q8 shiftv(x' 08 ',, x1 ,, ls ,, x1 )
call q8 xor v(x' 08 ',, x1 ,, dsp1 ,, dsp1)
x1=0
call vranf(ran,nh)
where (ran(1:nh).ge..5)x1=1
call q8 shiftv(x' 08 ',, x1 ,, ls ,, x1 )
call q8 xor v(x' 08 ',, x1 ,, dsp2 ,, dsp2)
1 continue

else if(init.eq.6)then

read(6,*)lo,betao,nso
do 76 i=1,ns
76 read(6,177) spin(i)
177 format(z16)

endif
return
end

subroutine raninit(xrand,np)
integer xrand(*)

c initializes the shift register random number generator.

do 1 i=1,np
ic=0
do 2 j=1,55
ic=shift(ic,1)
if((j.le.23).or.(j.ge.33.and.j.le.55))then
if(ranf().ge..5)ic=or(ic,1)
endif
2 continue
1 xrand(i)=ic
return
end

subroutine init64( i64 , n64 )

c sets i64 with integers in such a way as to satisfy that the serie
c [ i64(1) , i64(1)+i64(2) , ... , i64(1)+i64(2)+...+i64(64) ; (modulo 64) ]
c is a random permutation of the numbers 0 to 63.

dimension i64(64),n64(64)

n64(2:63)=0
i64(1)=0
n64(1)=1
ns=0
do 7 i=2,64
n=int(63.*ranf()+1)
nso=mod(n+ns,64)
if(n64(nso+1).eq.1)goto 7
ns=nso
7 n64(nso+1)=1
1 i64(i)=n
return
end

```

Fig. 2 (continued)

THE OUTPUT

BETA = 0.22165000000000  
 INITIAL LATTICE DISORDERED  
 THERMALIZING DURING 25000 SWEEPS  
 TAKING AVERAGES AFTER EVERY 100 SWEEPS DURING 25000 SWEEPS

LAT 1	M = -0.0048340	M =	S = 0.3407013	X =	110.175884	XP =	556.703504	C =	13.820515	GR =	-1.404986
LAT 2	M = 0.0060040	M =	S = 0.3424333	X =	100.785294	XP =	551.952328	C =	11.544944	GR =	-1.439802
LAT 3	M = 0.3419873	M =	S = 0.3419873	X =	107.300749	XP =	554.408006	C =	11.355583	GR =	-1.436903
LAT 4	M = 0.0091760	M =	S = 0.3406607	X =	96.555295	XP =	518.321312	C =	10.927884	GR =	-1.448633
LAT 5	M = 0.0095580	M =	S = 0.3410187	X =	97.222719	XP =	545.671716	C =	12.959921	GR =	-1.437544
LAT 6	M = 0.0282970	M =	S = 0.3386853	X =	88.297386	XP =	524.343018	C =	14.872295	GR =	-1.339440
LAT 7	M = -0.0178570	M =	S = 0.3392287	X =	106.889693	XP =	527.683138	C =	11.822033	GR =	-1.468217
LAT 8	M = 0.0533300	M =	S = 0.3417460	X =	103.688307	XP =	563.654308	C =	13.999393	GR =	-1.479494
LAT 9	M = 0.0011660	M =	S = 0.3429097	X =	94.774681	XP =	566.599204	C =	12.160427	GR =	-1.477912
LAT 10	M = -0.0645830	M =	S = 0.3415207	X =	105.472886	XP =	530.101158	C =	13.831400	GR =	-1.449696
LAT 11	M = -0.0644610	M =	S = 0.3399600	X =	107.025518	XP =	534.812118	C =	12.742202	GR =	-1.410785
LAT 12	M = 0.0571420	M =	S = 0.3432587	X =	102.934561	XP =	582.764092	C =	13.806341	GR =	-1.459483
LAT 13	M = 0.0178130	M =	S = 0.3471120	X =	101.954661	XP =	597.054638	C =	12.487388	GR =	-1.467532
LAT 14	M = 0.0222410	M =	S = 0.3416220	X =	105.205192	XP =	576.973202	C =	14.633021	GR =	-1.401232
LAT 15	M = 0.0511210	M =	S = 0.3448967	X =	110.701104	XP =	599.702266	C =	13.894701	GR =	-1.450949
LAT 16	M = -0.0228950	M =	S = 0.3416553	X =	101.638990	XP =	547.875188	C =	13.583701	GR =	-1.449402
LAT 17	M = 0.0063080	M =	S = 0.3404857	X =	99.561595	XP =	547.875188	C =	12.918553	GR =	-1.441420
LAT 18	M = -0.0149420	M =	S = 0.3484773	X =	99.608621	XP =	539.368248	C =	13.586372	GR =	-1.412280
LAT 19	M = 0.0048000	M =	S = 0.3421447	X =	100.684639	XP =	551.638488	C =	13.391780	GR =	-1.439817
LAT 20	M = 0.0653720	M =	S = 0.3426327	X =	102.462225	XP =	584.255064	C =	13.413788	GR =	-1.488519
LAT 21	M = -0.0358740	M =	S = 0.3414160	X =	103.887932	XP =	570.473660	C =	14.753841	GR =	-1.429721
LAT 22	M = -0.0398680	M =	S = 0.3415833	X =	99.504244	XP =	520.286268	C =	11.232732	GR =	-1.395200
LAT 23	M = 0.0172160	M =	S = 0.3415793	X =	107.736039	XP =	541.765516	C =	11.589478	GR =	-1.460724
LAT 24	M = 0.0144960	M =	S = 0.3397087	X =	108.871908	XP =	514.592228	C =	13.764910	GR =	-1.327357
LAT 25	M = 0.0042670	M =	S = 0.3397997	X =	105.065110	XP =	613.045022	C =	13.711048	GR =	-1.349436
LAT 26	M = 0.0410510	M =	S = 0.3440087	X =	112.043498	XP =	613.134086	C =	16.270822	GR =	-1.447704
LAT 27	M = -0.0317170	M =	S = 0.3397793	X =	109.410365	XP =	541.209610	C =	12.744961	GR =	-1.362801
LAT 28	M = -0.0167460	M =	S = 0.3232500	X =	101.122475	XP =	528.897976	C =	12.616406	GR =	-1.428730
LAT 29	M = -0.0065300	M =	S = 0.3391940	X =	100.711656	XP =	593.4231984	C =	11.692486	GR =	-1.356412
LAT 30	M = -0.0068300	M =	S = 0.3397800	X =	110.964414	XP =	523.492262	C =	12.241286	GR =	-1.355379
LAT 31	M = 0.0327330	M =	S = 0.3376960	X =	102.312602	XP =	596.395122	C =	14.611589	GR =	-1.317529
LAT 32	M = 0.0452070	M =	S = 0.3420780	X =	100.389207	XP =	617.351798	C =	13.795892	GR =	-1.523595
LAT 33	M = -0.0017660	M =	S = 0.3416613	X =	101.110378	XP =	564.171068	C =	12.906091	GR =	-1.448031
LAT 34	M = 0.0330590	M =	S = 0.3432627	X =	100.807998	XP =	586.633706	C =	14.828275	GR =	-1.449726
LAT 35	M = 0.0024570	M =	S = 0.3415588	X =	106.043696	XP =	552.572358	C =	14.470135	GR =	-1.375563
LAT 36	M = -0.0095780	M =	S = 0.3402893	X =	106.734894	XP =	529.338944	C =	11.459671	GR =	-1.417201
LAT 37	M = -0.0225300	M =	S = 0.3401787	X =	94.498939	XP =	541.540796	C =	13.412455	GR =	-1.444748
LAT 38	M = -0.0225300	M =	S = 0.3375653	X =	103.354811	XP =	494.167680	C =	12.974650	GR =	-1.349329
LAT 39	M = 0.0009450	M =	S = 0.3440880	X =	106.708954	XP =	583.363110	C =	13.392845	GR =	-1.485523
LAT 40	M = -0.0502000	M =	S = 0.3413707	X =	105.732957	XP =	553.667868	C =	14.063455	GR =	-1.469034
LAT 41	M = 0.0248120	M =	S = 0.3408440	X =	104.527774	XP =	511.752568	C =	11.817144	GR =	-1.313046
LAT 42	M = 0.0417440	M =	S = 0.3393200	X =	103.064641	XP =	552.675658	C =	11.871444	GR =	-1.313046
LAT 43	M = 0.0138510	M =	S = 0.3417393	X =	103.064641	XP =	552.675658	C =	13.372450	GR =	-1.433000
LAT 44	M = 0.0119640	M =	S = 0.3352530	X =	107.540090	XP =	550.291052	C =	11.060873	GR =	-1.433784
LAT 45	M = -0.0391620	M =	S = 0.3307960	X =	102.290283	XP =	528.3424632	C =	13.147762	GR =	-1.372322
LAT 46	M = 0.0181780	M =	S = 0.3398800	X =	110.263588	XP =	516.185712	C =	14.517491	GR =	-1.330908
			S = 0.3403007	X =	99.446916	XP =	531.874596	C =	12.709390	GR =	-1.440305



LAT 47	M = 0.0080730	M = 0.2440050	S = 0.3418740	X = 93.091318	XP = 569.398838	C = 12.516944	GR = -1.485784
LAT 48	M = 0.0070040	M = 0.2426860	S = 0.3434460	X = 96.809479	XP = 567.981436	C = 12.178103	GR = -1.476262
LAT 49	M = 0.0457270	M = 0.3416933	S = 0.3416933	X = 100.120920	XP = 556.656910	C = 12.530724	GR = -1.451987
LAT 50	M = -0.0133020	M = 0.2383760	S = 0.3425980	X = 95.285625	XP = 549.870564	C = 13.907340	GR = -1.460578
LAT 51	M = 0.0433750	M = 0.3410790	S = 0.3410790	X = 95.596766	XP = 544.967726	C = 11.677170	GR = -1.480204
LAT 52	M = 0.0024720	M = 0.2460048	S = 0.3447127	X = 111.689312	XP = 595.803056	C = 12.118731	GR = -1.453982
LAT 53	M = 0.0310740	M = 0.2358400	S = 0.3407327	X = 99.688343	XP = 544.652386	C = 13.143439	GR = -1.403298
LAT 54	M = -0.0165520	M = 0.2375160	S = 0.3416146	X = 102.624162	XP = 553.934964	C = 14.600107	GR = -1.424157
LAT 55	M = 0.0033870	M = 0.2414870	S = 0.3404999	X = 93.884913	XP = 569.417782	C = 14.944336	GR = -1.443537
LAT 56	M = 0.0070730	M = 0.2311990	S = 0.3402113	X = 101.265161	XP = 528.889982	C = 11.991000	GR = -1.437309
LAT 57	M = 0.0424970	M = 0.2380810	S = 0.3413160	X = 103.730126	XP = 557.190626	C = 14.273246	GR = -1.413908
LAT 58	M = 0.0267990	M = 0.2459770	S = 0.3418933	X = 102.415334	XP = 586.452810	C = 13.440079	GR = -1.470905
LAT 59	M = 0.0057930	M = 0.2307130	S = 0.3408100	X = 110.073627	XP = 535.901534	C = 14.677312	GR = -1.364311
LAT 60	M = -0.0756560	M = 0.2382820	S = 0.3410297	X = 93.147560	XP = 547.374052	C = 12.300062	GR = -1.438487
LAT 61	M = 0.0334920	M = 0.2270960	S = 0.3399493	X = 98.100586	XP = 510.681332	C = 11.204285	GR = -1.407199
LAT 62	M = 0.0207010	M = 0.2370960	S = 0.3370800	X = 94.834708	XP = 500.039334	C = 12.393889	GR = -1.427377
LAT 63	M = -0.0381700	M = 0.2409560	S = 0.3420673	X = 110.739259	XP = 602.993492	C = 14.359734	GR = -1.471558
LAT 64	M = 0.0069490	M = 0.2320050	S = 0.3407167	X = 107.225510	XP = 537.836070	C = 11.977657	GR = -1.378928

DATA FROM AVERAGING THE 64 LATTICES :

M = 0.00394073	M  = 0.23622336	Mt2 = 0.06667325	Mt4 = 0.00744216
	+/- 0.00091810	+/- 0.00044161	+/- 0.00007913
	S = 0.34115696	S+2 = 0.11693744	S+4 = 0.01393827
	+/- 0.00019446	+/- 0.00013363	+/- 0.00003215
	X = 102.54262	C = 13.126715	GR = -1.4226765
	+/- 0.66730	+/- 0.144369	+/- 0.00598663

XP = 549.38599  
+/- 3.53290

EFFECTIVE BETA = 0.2216629017657

CORRECTED DATA:

M = 0.00394073	M  = 0.23554687	Mt2 = 0.06634626	Mt4 = 0.00738606
	+/- 0.00091810	+/- 0.00044161	+/- 0.00007913
	S = 0.34098685	S+2 = 0.11681966	S+4 = 0.01390967
	+/- 0.00019446	+/- 0.00013363	+/- 0.00003215
	X = 102.90392	C = 13.142999	GR = -1.4188134
	+/- 0.66730	+/- 0.144369	+/- 0.00598663

XP = 546.77009  
+/- 3.53290

TIME SPENT IN THE MAIN PARTS OF THE PROGRAM :

PROGRAM ISING =	288.27 SECONDS
IN UPDATING =	266.21 SECONDS
IN MEASURING =	27.28 SECONDS
SUB. DEMETRO =	15.73 SECONDS
SUB. INITNR =	3.92 SECONDS
SUB. INIT64 =	1.01 SECONDS
SUB. MEAS =	27.25 SECONDS

RUNNING AT 98.4 MEGAFLIPS

Fig. 3. Output obtained from running the code shown in Fig. 2. The parameters used are those given in the code.  $M$  stands for the magnetization per spin  $\langle M \rangle / V$ ,  $XP$  for  $\langle M^2 \rangle / V$ ,  $X$  for  $\langle M^2 \rangle - \langle M \rangle^2 / V$ ,  $GR$  for the renormalized coupling constant  $(M^4) / \langle M^2 \rangle^2 - 3$ ,  $S$  for the nearest-neighbor spin-spin correlation function  $\langle s_i s_{i+h} \rangle$ , and  $C$  for  $dS/d\beta$ . The data is corrected using eq. 10.

distribution. The corrected data is shown in Fig. 1(b). Note that each data point is corrected by a different amount because each corresponds to a different stream of random numbers and therefore, a different  $\beta_{\text{eff}}$ . In Fig. 1(c) we plot a conventional Metropolis run done by using completely independent sets of random numbers for the 64 lattices. The vertical line corresponds to the exact value. We have repeated this experiment three more times for a total of 160 runs. In 106 and 103 of these runs,  $\langle S \rangle$  and  $\langle S^2 \rangle$  were within 1 sd of the exact value. This corresponds to 66.25% and 64.37%, respectively, and indicates that the errors are properly defined.

## THE CODE

The code is written in standard Cyber Fortran 200 using the special Q8 calls which translate directly into machine code. We use descriptors to point to arrays in the standard way, and the motivated reader is directed to the Fortran 200 manual for inspiration. A listing of the code is included with this paper along with the output (Fig. 2 and 3).<sup>3</sup>

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<sup>3</sup> Fortran 200 Version 1 Reference Manual (publication number 60485000) CDC Cyber 205 Hardware Reference Manual (publication number 60256020). These manuals are available from: Control Data Corporation, Literature and Distribution Services, 308 North Dale Street, St. Paul, MN 55103.

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